

EECS3311 Software Design (Fall 2020)

Q&A - Lecture Series W10

Monday, November 23

Subcontracting: Architectural View

$p \Rightarrow q$

harder to satisfy?



my_phone



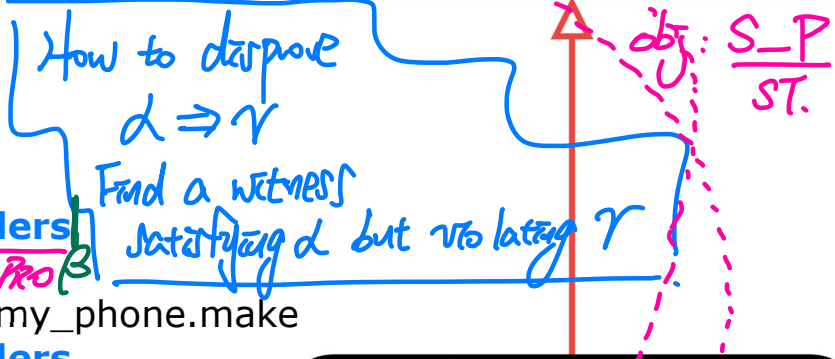
my_phone: SMART_PHONE

create my_phone.make

list := my_phone.get_reminders

create {IPHONE_11_PRO} my_phone.make

list := my_phone.get_reminders

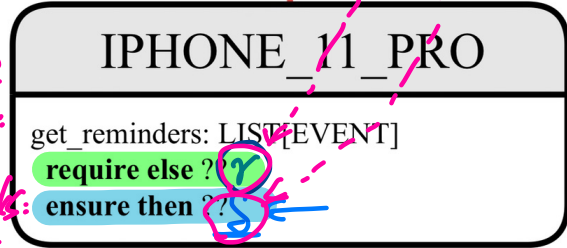


(design time) \rightarrow paper
 child precond requires less:
 weaker

(runtime) \rightarrow execute
 child precondition checks:
 $P: [d] \vee [\gamma] C.$

$d \Rightarrow \gamma$
 child postcond
 $S \Rightarrow B.$
 EASIER TO SATISFY
 ENFORCE MORE
 STRONGER

child postcondition checks:
 $P: [B] \wedge [S] C.$
 HARDER TO SATISFY



Subcontracting: Example (1)

```
class SMART_PHONE
  get_reminders: LIST[EVENT]
  require
     $\alpha$ : battery_level  $\geq$  0.1 -- 10%
  ensure
     $\beta$ :  $\forall e$ :Result | e happens today
end
```

satisfies

12%

```
class IPHONE_11_PRO
  inherit SMART_PHONE redefine get_reminders end
  get_reminders: LIST[EVENT]
  require else
     $\gamma$ : battery_level  $\geq$  0.15 -- 15%
  ensure then
     $\delta$ :  $\forall e$ :Result | e happens today or tomorrow
  end
```

violates

stronger

requires more

poor design.

The computer does not
prove or disprove:
 $d \Rightarrow \gamma$

At runtime:

$T \circ d \vee \gamma \circ F = T$

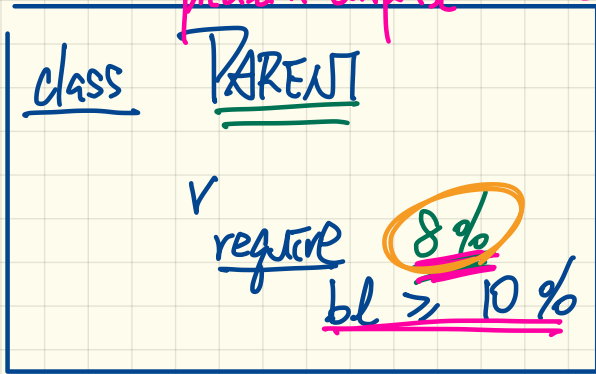
```
my_phone: SMART_PHONE
```

```
create my_phone.make
list := my_phone.get_reminders
```

```
create {IPHONE_11_PRO} mine.make
list := my_phone.get_reminders
```



shock: what works in parent does not work in child.
 pleasant surprise: what fails to work in parent actually works in child.

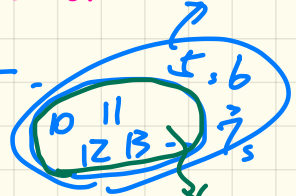


Design Time

True

$bl \geq 10\% \Rightarrow$

$bl \geq 5\%$
weaker.

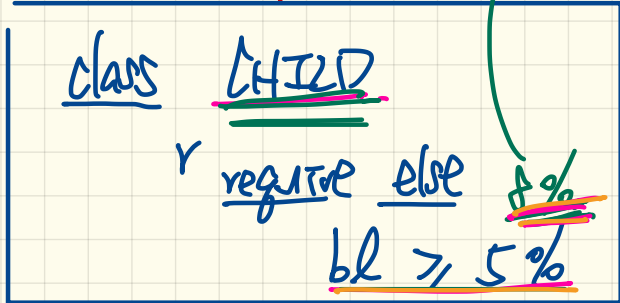
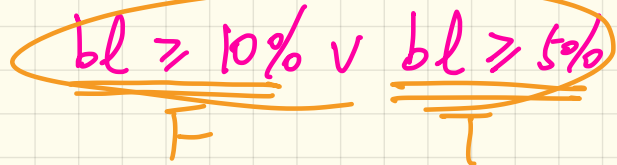


more accommodating for the input value

Runtime obj. r

if DT of obj is CHILD

check: I (no precond violation)



Subcontracting: Example (2)

```

class SMART_PHONE
  get_reminders: LIST[EVENT]
  require
    α: battery_level ≥ 0.1 -- 10%
  ensure
    β: ∀e:Result | e happens today
end
  
```

Imp: Return a list of (time) events.

violates (arrow from β to Imp)

```

class IPHONE_11_PRO
  inherit SMART_PHONE redefine get_reminders end
  get_reminders: LIST[EVENT]
  require else
    γ: battery_level ≥ 0.15 -- 15%
  ensure then
    δ: ∀e:Result | e happens today or tomorrow
  end
  
```

violates (arrow from δ to Imp)

Design Time

Want to prove:

$$\underline{\underline{\delta}} \Rightarrow \underline{\underline{\beta}} \quad F \quad T \quad \boxed{\beta} \wedge \boxed{\delta} = \boxed{F}$$

postcond. violation (arrow from F to δ)

Runtime:

check for DT IP-11 Pro:

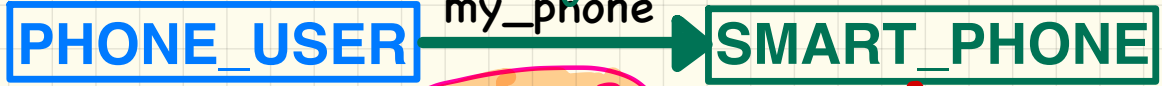
```

my_phone: SMART_PHONE

create my_phone.make
list := my_phone.get_reminders

create {IPHONE_11_PRO} mine.make
list := my_phone.get_reminders
  
```

$\beta \Rightarrow \delta$ fact \Rightarrow poor design



β today δ tomorrow



How is a Poor Design Checked at Runtime?

```
class FOO
  f
  require
    original_pre  $\alpha$ 
  ensure
    original_post  $\beta$ 
end
```

```
class BAR
  inherit FOO redefine f end
  f
  require else
    new_pre  $\gamma$ 
  ensure then
    new_post  $\delta$ 
  end
end
```

(Static) Design Time :

- $original_pre \Rightarrow new_pre$ should be proved as a tautology
- $new_post \Rightarrow original_post$ should be proved as a tautology

(Dynamic) Runtime :

- $original_pre \vee new_pre$ is checked
- $original_post \wedge new_post$ is checked

Even if these two logical implications are not implemented during compile time due to poor design, the compiler won't complain and it does not lead to a run time violation but it may only give shock to the user? *not accurate.*

① $orig_pre \not\Rightarrow new_pre$
↓ as long as $orig_pre$ is satisfied, new_pre is not.
no precond. violation

② $new_post \not\Rightarrow orig_post$
↓ as long as $orig_post$ is violated, there's postcond. violation.

In practice

PARENT

γ
require
 α
ensure

β

CHILD

redefine γ end

① γ do end \equiv ② γ

(not equivalent)

Is this a good design? γ

require else $\boxed{\text{True}}$

$d \Rightarrow \text{True}$
 $\equiv \underline{\underline{\text{True}}}$

RT: $\alpha \vee \text{True}$
 $\equiv \underline{\underline{\text{True}}}$

no input restriction
at all

RT: $\beta \vee \text{False}$
 $\equiv \underline{\underline{\beta}}$

RT: $\beta \wedge \text{True}$
 $\equiv \underline{\underline{\beta}}$

Multiple Inheritance: Exercise

```
class RECTANGLE
  feature -- Queries
    width, height: REAL
    xpos, ypos: REAL
  feature -- Commands
    make (w, h: REAL)
    change_width
    change_height
    move
end
```

```
class TREE[G]
  feature -- Queries
    descendants: ITERABLE[G]
  feature -- Commands
    add (c: G)
      -- Add a child 'c'.
end
```

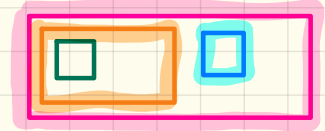
children at top level

```
class WINDOW
  inherit
    RECTANGLE
    TREE[WINDOW]
end
```

```
test_window: BOOLEAN
local w1, w2, w3, w4: WINDOW
do
  create w1 make(8, 6) ; create w2.make(4, 3)
  create w3.make(1, 1) ; create w4.make(1, 1)
  w2.add(w4) ; w1.add(w2) ; w1.add(w3)
  Result := w1.descendants.count = 2
end
```

You may:

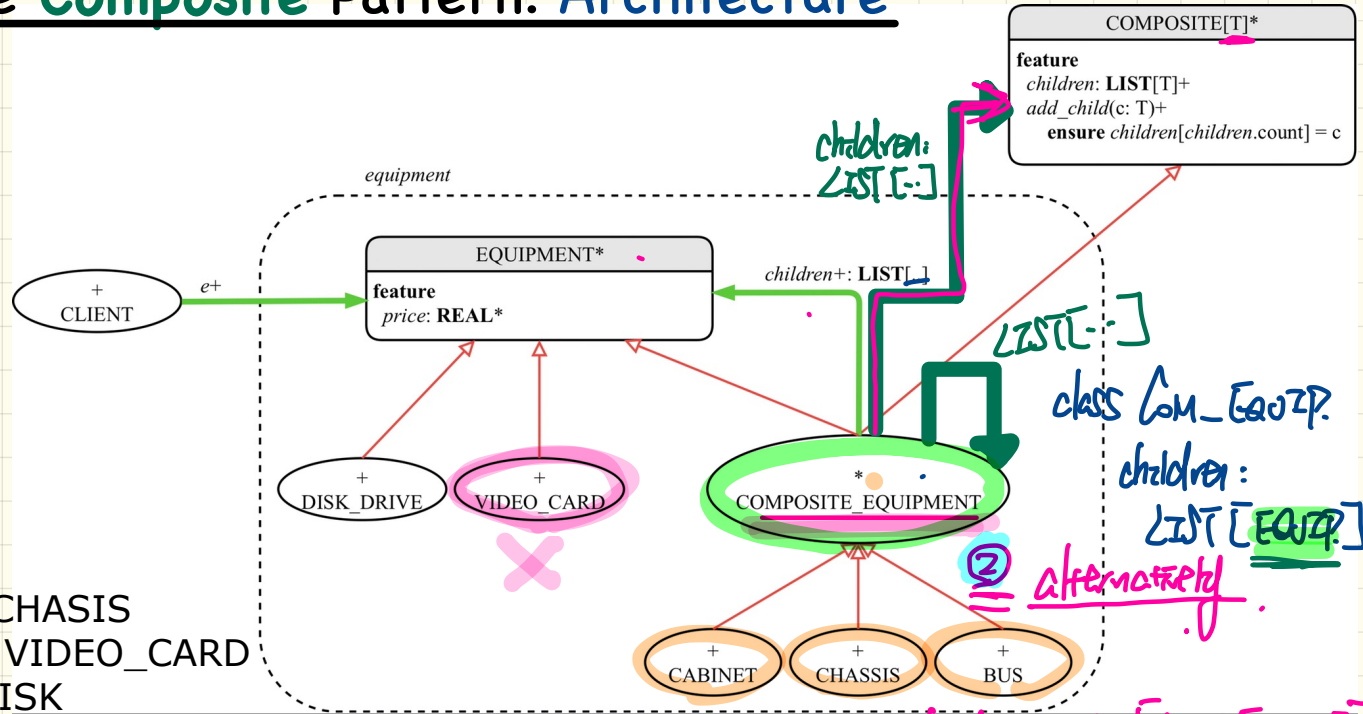
- ① rename 'descendants' to 'children'
- ② implement descendants to



include all nested window

Why does `w1.descendants` have 2 items instead of 4 (`w1`, `w2`, `w3`, `w4`)?

The Composite Pattern: Architecture



ch: CHASSIS
 crd: VIDEO_CARD
 d: DISK

create ch.make
create crd.make
create d.make

ch.add_child(crd) ✗
 ch.add_child(d) ✗
 crd.add_child(d) ✗

Is `children: LIST[COMPOSITE]` a valid design?

Invalid
 ∴ we cannot add any basic equip to the list.

children: LIST[COMP.[EQUIP.]

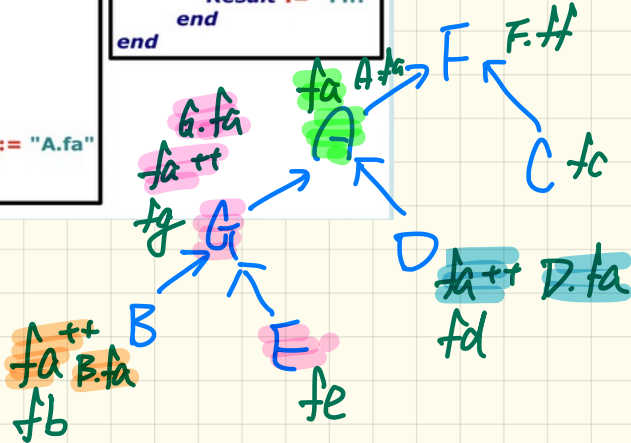
children: LIST[COM..EQU]

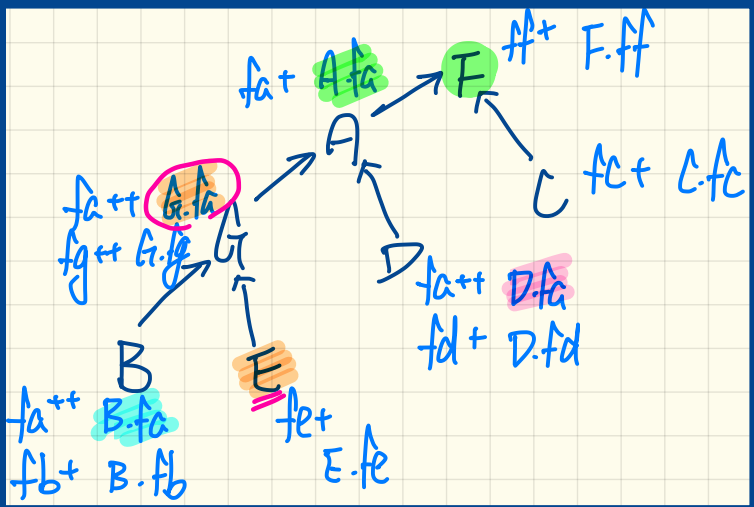
class Com_Equip.
 children:
 LIST[EQUIP.]
 alternately.

Quiz8: Polymorphic Return Values

<pre>class B inherit G redefine fa end feature fa: STRING do Result := "B.fa" end fb: STRING do Result := "B.fb" end end</pre>	<pre>class D inherit A redefine fa end feature fa: STRING do Result := "D.fa" end fd: STRING do Result := "D.fd" end end</pre>	<pre>class G inherit A redefine fa end feature fa: STRING do Result := "G.fa" end fg: STRING do Result := "G.fg" end end</pre>
<pre>class C inherit F feature fc: STRING do Result := "C.fc" end end</pre>	<pre>class E inherit G feature fe: STRING do Result := "E.fe" end end</pre>	<pre>class A inherit F feature fa: STRING do Result := "A.fa" end end</pre>

```
class F
feature
  ff: STRING
  do
    Result := "F.ff"
  end
end
```



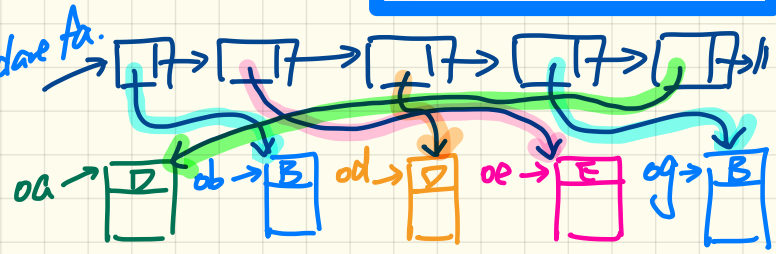


```

class
  CLIENT_2
create
  make
feature -- Collection
  a: LIST[A]
  make
  do
    create {LINKED_LIST[A]} a.make
  end
feature -- Routines
  ✓ add_a (obj: A)
  do
    a.extend (obj)
  end
  ✓ get_a (i: INTEGER): A
  require
    1 <= i and i <= a.count
  do
    ✓ st: A
    ✓ st: A
    Result := a[i]
  end
  ✓ get_f (i: INTEGER): F
  require
    1 <= i and i <= a.count
  do
    ✓ st: F
    ✓ st: A
    Result := a[i]
  end
end
  
```

oa: A
ob: B
od: D
oe: E
of: F
og: G
 client: CLIENT_2
 create {D} oa.make
 create {B} ob.make
 create {D} od.make
 create {E} oe.make
 create {B} og.make
 create client.make
 client.add_a (ob)
 client.add_a (oe)
 client.add_a (od)
 client.add_a (og)
 client.add_a (oa)

client.get_a(2).fa	Returns "G.fa" ⇩
client.get_f(2).fa	Invalid ⇩
client.get_a(5).fa	Returns "D.fa" ⇩
client.get_a(4).fa	Returns "B.fa" ⇩
client.get_a(1).fa	Returns "B.fa" ⇩
od := client.get_a(5)	Invalid ⇩
client.get_a(3).fa	Returns "D.fa" ⇩
og := client.get_a(1)	Invalid ⇩
client.get_f(4).fa	Invalid ⇩
of := client.get_a(4)	Valid ⇩



polymorphic collection